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# *Methods for Intelligent Systems*

## *Lecture Notes on Machine Learning*

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## *Probability for Dataminers*

### *– Probability Basics –*

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## Probability and Boolean Random Variables

**Boolean-valued random variable**  $A$  is a Boolean-valued random variable if  $A$  denotes an event, and there is some degree of uncertainty as to whether  $A$  occurs.

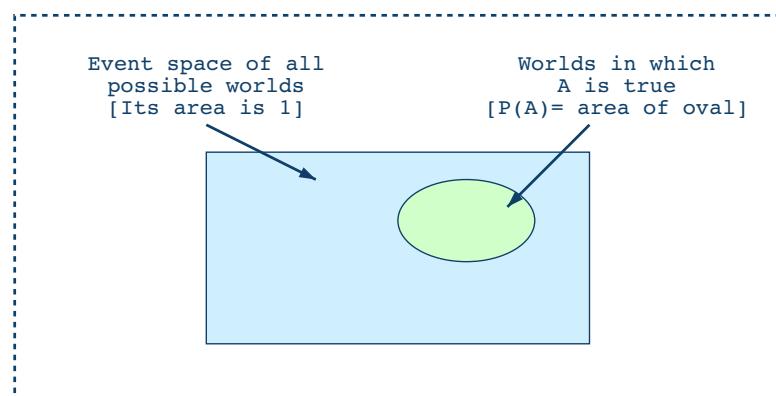
- Examples
  - $A =$  The US president in 2023 will be male
  - $A =$  You wake up tomorrow with a headache
  - $A =$  You like the “Gladiator”

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## Probability and Boolean Random Variables

**Boolean-valued random variable**  $A$  is a Boolean-valued random variable if  $A$  denotes an event, and there is some degree of uncertainty as to whether  $A$  occurs.

**Probability of  $A$**  “the fraction of possible worlds in which  $A$  is true”



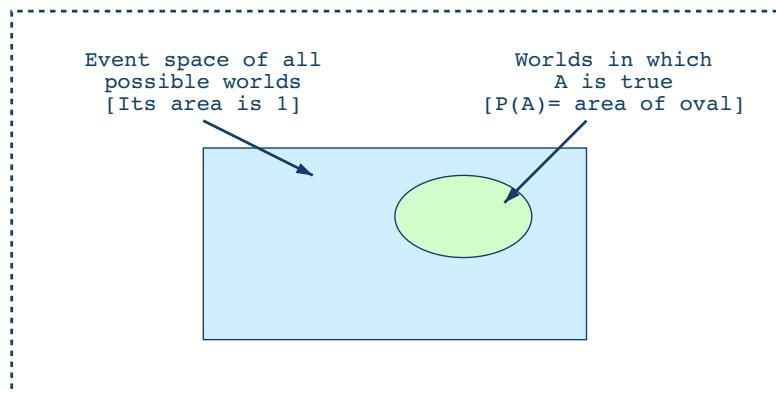
Note: this is one of the possible definitions. We won't go into the philosophy of it!

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## Probability Axioms

Define the whole set of possible worlds with the label `true` and the empty set with `false`:

- $0 \leq P(A) \leq 1$
- $P(A = \text{true}) = 1; P(A = \text{false}) = 0$
- $P(A \vee B) = P(A) + P(B) - P(A \wedge B)$

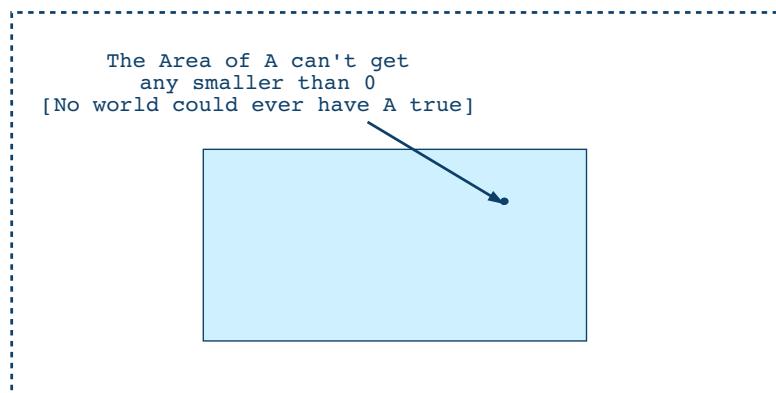


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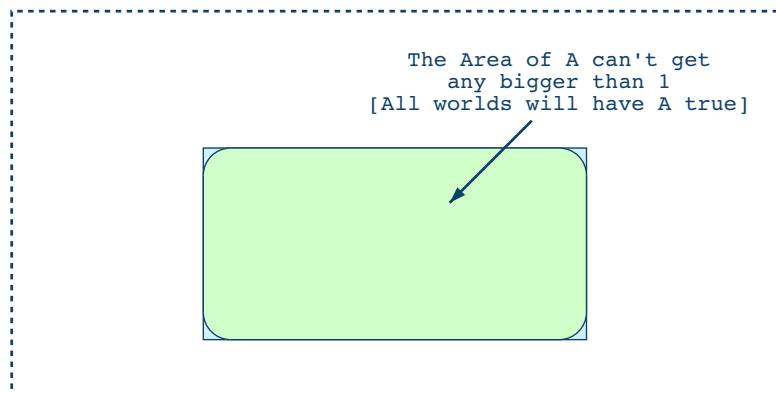


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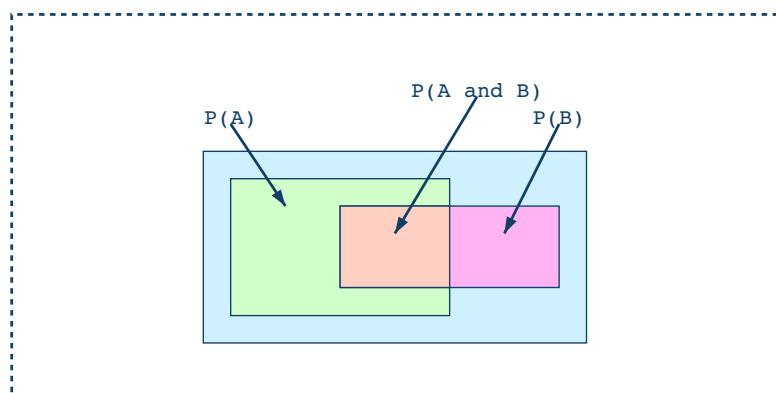


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## Probability Axioms

Define the whole set of possible worlds with the label `true` and the empty set with `false`:

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## Theorems From the Axioms (I)

Using the axioms:

- $P(A = \text{true}) = 1; P(A = \text{false}) = 0$
- $P(A \vee B) = P(A) + P(B) - P(A \wedge B)$

Proove:  $P(\sim A) = P(\bar{A}) = 1 - P(A)$

$$\begin{aligned}\text{true} &= A \vee \bar{A} \\ P(\text{true}) &= P(A \vee \bar{A}) \\ &= P(A) + P(\bar{A}) - P(A \wedge \bar{A}) \\ &= P(A) + P(\bar{A}) - P(\text{false}) \\ 1 &= P(A) + P(\bar{A}) - 0 \\ 1 - P(A) &= P(\bar{A})\end{aligned}$$

## Theorems From the Axioms (II)

Using the axioms:

- $P(A = \text{true}) = 1; P(A = \text{false}) = 0$
- $P(A \vee B) = P(A) + P(B) - P(A \wedge B)$

Proove:  $P(A) = P(A \wedge B) + P(A \wedge \bar{B})$

$$\begin{aligned}A &= A \wedge \text{true} \\ &= A \wedge (B \vee \bar{B}) \\ &= (A \wedge B) \vee (A \wedge \bar{B}) \\ P(A) &= P((A \wedge B) \vee (A \wedge \bar{B})) \\ &= P(A \wedge B) + P(A \wedge \bar{B}) - P((A \wedge B) \wedge (A \wedge \bar{B})) \\ &= P(A \wedge B) + P(A \wedge \bar{B}) - P(\text{false}) \\ &= P(A \wedge B) + P(A \wedge \bar{B})\end{aligned}$$

## Multivalued Random Variables

**Multivalued random variable**  $A$  is a *random variable of arity  $k$*  if it can take on exactly one values out of  $\{v_1, v_2, \dots, v_k\}$ .

We still have the probability axioms plus

- $P(A = v_i \wedge A = v_j) = 0$  if  $i \neq j$
- $P(A = v_1 \vee A = v_2 \vee \dots \vee A = v_k) = 1$

Proove:  $P(A = v_1 \vee A = v_2 \vee \dots \vee A = v_i) = \sum_{j=1}^i P(A = v_j)$

Proove:  $\sum_{j=1}^k P(A = v_j) = 1$

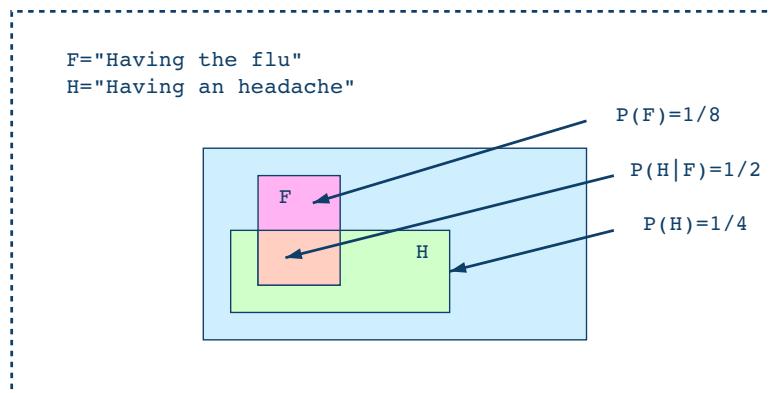
Proove:

$P(B \wedge [A = v_1 \vee A = v_2 \vee \dots \vee A = v_i]) = \sum_{j=1}^i P(B \wedge A = v_j)$

Proove:  $P(B) = \sum_{j=1}^k P(B \wedge A = v_j)$

## Conditional Probability

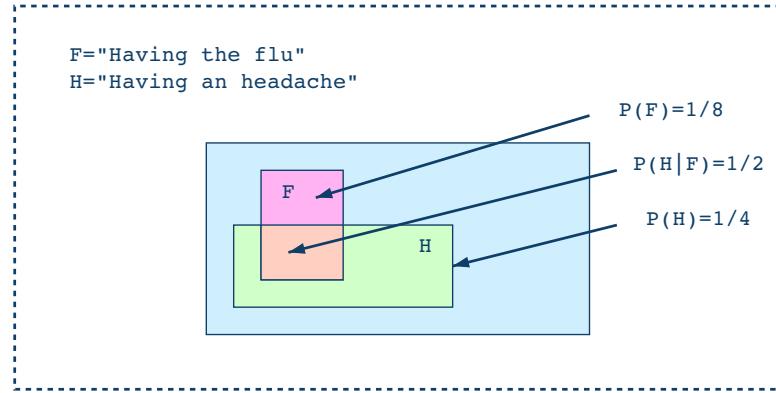
**Probability of  $A$  given  $B$ :** “the fraction of possible worlds in which  $B$  is true that also have  $A$  true”



“Sometimes I’ve the flu and sometimes I’ve a headache, but half of the times I’m with the flu I’ve also a headache!”

## Conditional Probability

**Probability of  $A$  given  $B$ :** “the fraction of possible worlds in which  $B$  is true that also have  $A$  true”

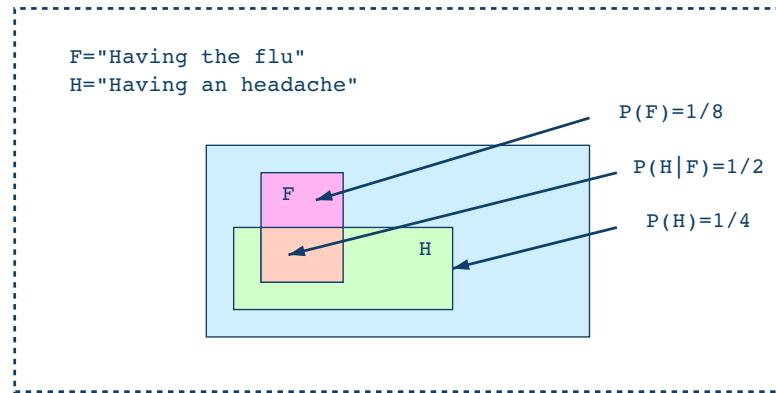


$$P(H|F) = \frac{\text{Num. of worlds with } F \text{ and } H}{\text{Num. worlds with } F} = \frac{P(H \wedge F)}{P(F)}$$

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## Probabilistic Inference

One day you wake up with a headache and you think: “*Half of the flus are associated with headaches so I must have 50% chance of getting the flu*”.



Is this reasoning correct?

$$P(F|H) = \frac{P(F \wedge H)}{P(H)} = \frac{P(H \wedge F)}{P(H)} = \frac{P(H|F) * P(F)}{P(H)} = \frac{1/2 * 1/8}{1/4} = 1/4$$

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## Theorems that we used (and will use)

In doing the previous inference we have used two famous theorems:

- Chain rule

$$P(A \wedge B) = P(A|B)P(B)$$

- Bayes theorem

$$P(A|B) = \frac{P(A \wedge B)}{P(B)} = \frac{P(B|A)P(A)}{P(B)}$$

We can have more general formulae:

- $P(A|B) = \frac{P(B|A)P(A)}{P(B|A)P(A) + P(B|\bar{A})P(\bar{A})}$
- $P(A|B \wedge X) = \frac{P(B|A \wedge X)P(A \wedge X)}{P(B \wedge X)}$
- $P(A = v_i|B) = \frac{P(B|A=v_i)P(A=v_i)}{\sum_{k=1}^{n_A} P(B|A=v_k)P(A=v_k)}$

## Independent Variables

**Independent variables:** Assume  $A$  and  $B$  are boolean random variables;  $A$  and  $B$  are independent (denote it with  $A \perp B$ ) if and only if:

$$P(A|B) = P(A)$$

Using the definition:

- $P(A|B) = P(A)$

Proove:  $P(A \wedge B) = P(A)P(B)$

$$\begin{aligned} P(A \wedge B) &= P(A|B)P(B) \\ &= P(A)P(B) \end{aligned}$$

Proove:  $P(B|A) = P(B)$

$$P(B|A) = \frac{P(A|B)P(B)}{P(A)}$$

# Probability for Dataminers

## – Information Gain –

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### Information and Bits

Your mission, if you decide to accept it, will be:

*“Transmit a set of independent random samples of  $X$  over a binary serial link.”*

1. Starring at  $X$  for a while, you notice that it has only four possible values: A, B, C, D
2. You decide to transmit the data encoding each reading with two bits:

$$A = 00, B = 01, C = 10, D = 11.$$

Mission Accomplished!

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## Information and “Fewer Bits”

Your mission, if you decide to accept it, will be:

“The previous code uses 2 bits for symbol.

Knowing that the probabilities are not equal:  $P(X=A)=1/2$ ,  $P(X=B)=1/4$ ,  $P(X=C)=1/8$ ,  $P(X=D)=1/8$ , invent a coding for your transmission that only uses 1.75 bits on average per symbol.”

1. You decide to transmit the data encoding each reading with a different number of bits:

$$A = 0, B = 10, C = 110, D = 111.$$

Mission Accomplished!

## Information and Entropy

Suppose  $X$  can have one of  $m$  values with probability

$$P(X = V_1) = p_1, \dots, P(X = V_m) = p_m.$$

What's the smallest possible number of bits, on average, per symbol, needed to transmit a stream of symbols drawn from  $X$ 's distribution?

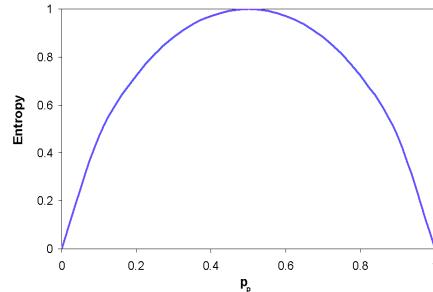
$$\begin{aligned} H(X) &= -p_1 \log_2 p_1 - p_2 \log_2 p_2 - \dots - p_m \log_2 p_m \\ &= -\sum_{j=1}^m p_j \log_2 p_j = \text{Entropy of } X \end{aligned}$$

“Good idea! But what is entropy anyway?”

## Entropy: “What is it anyway?”

### Simple Case:

- $X$  has 2 values  $\oplus$  and  $\ominus$
- $p_{\oplus}$  probability of  $\oplus$
- $p_{\ominus} = 1 - p_{\oplus}$  probability of  $\ominus$



$$H(X) = -p_{\ominus} \log_2 p_{\ominus} - p_{\oplus} \log_2 p_{\oplus}$$

Entropy measures “disorder” or “uniformity in distribution”

1. *High Entropy*:  $X$  is very “disordered”  $\rightarrow$  “interesting”
2. *Low Entropy*:  $X$  is very “ordered”  $\rightarrow$  “boring”

## Useful Facts on Logarithms

Just for you to know it might be useful to review a couple of formulas to be used in calculation:

- $\ln x \times y = \ln x + \ln y$
- $\ln \frac{x}{y} = \ln x - \ln y$
- $\ln x^y = y \times \ln x$
- $\log_2 x = \frac{\ln x}{\ln 2} = \frac{\log_{10} x}{\log_{10} 2}$
- $\log_a x = \frac{\ln x}{\ln a}$
- $\log_2 0 = -\infty$  (the formula is no good for a probability of 0)

Now we can practice with a simple example!

## Specific Conditional Entropy

Suppose we are interested in predicting output  $Y$  from input  $X$  where

| X       | Y   |
|---------|-----|
| Math    | Yes |
| History | No  |
| CS      | Yes |
| Math    | No  |
| Math    | No  |
| CS      | Yes |
| History | No  |
| Math    | Yes |

- $X$  = University subject
- $Y$  = Likes the movie “Gladiator”

From this data we can estimate

- $P(Y = \text{Yes}) = 0.5$
- $P(X = \text{Math}) = 0.5$
- $P(Y = \text{Yes} | X = \text{History}) = 0$

Definition of Specific Conditional Entropy:

- $H(Y|X=v)$ : *the entropy of  $Y$  only for those records in which  $X$  has value  $v$* 
  - $H(Y|X=\text{Math}) = 1$
  - $H(Y|X=\text{History}) = 0$

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## Conditional Entropy

Definition of Conditional Entropy  $H(Y|X)$ :

- *The average  $Y$  specific conditional entropy*
- *Expected number of bits to transmit  $Y$  if both sides will know the value of  $X$*
- $\sum_j P(X = v_j)H(Y|X = v_j)$

Definition of Conditional Entropy  $H(Y|X)$ :

- $\sum_j P(X = v_j)H(Y|X = v_j)$

| X       | Y   |
|---------|-----|
| Math    | Yes |
| History | No  |
| CS      | Yes |
| Math    | No  |
| Math    | No  |
| CS      | Yes |
| History | No  |
| Math    | Yes |

| $v_j$   | $P(X = v_j)$ | $H(Y X = v_j)$ |
|---------|--------------|----------------|
| Math    | 0.5          | 1              |
| History | 0.25         | 0              |
| CS      | 0.25         | 0              |

$$H(Y|X) = ?$$

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## Information Gain

| X       | Y   |
|---------|-----|
| Math    | Yes |
| History | No  |
| CS      | Yes |
| Math    | No  |
| Math    | No  |
| CS      | Yes |
| Hystory | No  |
| Math    | Yes |

*I must transmit Y on a binary serial line.  
How many bits on average would it save me if both  
ends of the line knew X?*

$$\begin{aligned}IG(Y|X) &= H(Y) - H(Y|X) \\ &= 1 - 0.5 = 0.5\end{aligned}$$

Information Gain measures the “information”  
provided by  $X$  to predict  $Y$

This IS about Machine Learning!

## Relative Information Gain

| X       | Y   |
|---------|-----|
| Math    | Yes |
| History | No  |
| CS      | Yes |
| Math    | No  |
| Math    | No  |
| CS      | Yes |
| Hystory | No  |
| Math    | Yes |

*I must transmit Y on a binary serial line.  
What fraction of the bits on average would it save  
me if both ends of the line knew X?*

$$\begin{aligned}RIG(Y|X) &= (H(Y) - H(Y|X))/H(Y) \\ &= (1 - 0.5)/1 = 0.5\end{aligned}$$

Well, we'll find soon Information Gain and Relative  
Information gain talking about supervised learning  
with Decision Trees ...

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## Why is Information Gain Useful?

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Your mission, if you decide to accept it, will be:

*“Predict whether someone is going live  
past 80 years.”*

From historical data you might find:

- $IG(\text{LongLife} \mid \text{HairColor}) = 0.01$
- $IG(\text{LongLife} \mid \text{Smoker}) = 0.2$
- $IG(\text{LongLife} \mid \text{Gender}) = 0.25$
- $IG(\text{LongLife} \mid \text{LastDigitOfSSN}) = 0.00001$

What you should look at?